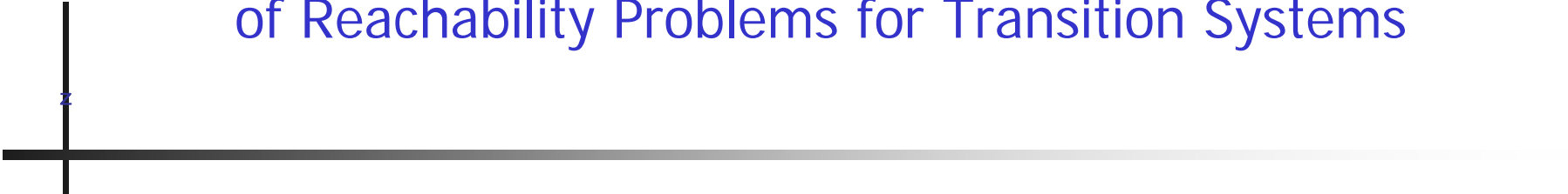


Why Multi-Result Supercompilation Matters: Case Study of Reachability Problems for Transition Systems



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July 8, 2012

Third International Valentin Turchin Workshop on Metacomputation, META 2010
Pereslavl-Zalessky, Russia

History and Main Conclusion

- **2005-2007 Andrei Nemytykh and Alexei Lisitsa**
 - have **experimentally** found a method to solve the **coverability problem** for (a class of) **practical counter systems** (models of cache-coherence protocols and other systems) with the Refal Supercompiler SCP4
 - *User + single-result supercompiler = MRSC*
- **2010-2011 Andrei Klimov**
 - have **theoretically** explained and proved that the **coverability problem** is solvable for **monotonic counter systems** by an **iterative procedure** of applying a **domain-specific supercompiler** for counter systems **varying a parameter** of generalization
 - *An optimized MRSC enumerating a small subset of residual graphs*
- **2005-2007 Gilles Geeraerts *et al* (Belgium)**
 - theory of 'Expand, Enlarge and Check' algorithmic schema (ECC) for solving the **coverability problem** of **well-structured transition systems** (WSTS)
 - *An MRSC for WSTS with reduced search space*

One thing to remember from this talk

- These are instances of **domain-specific multi-result supercompilation (MRSC)** with **search space reduction** based of **domain properties and purpose**

`User-controlled' MRSC

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 - have **experimentally** found a method to solve the **coverability problem** for (a class of) **practical counter systems** (models of cache-coherence protocols and other systems) with the Refal Supercompiler SCP4
 - *User + single-result supercompiler = MRSC*
- **Andrei Nemytykh devised two versions of SCP4**
 - SCP4₀ – standard version
 - SCP4₁ – generalization of empty expressions (representing zeros) prohibited
- **The user behavior**
 - When SCP4₀ did not prove the coverability, SCP4₁ was applied
 - No more supercompilers were needed for the considered samples borrowed from the collection by Giorgio Delzanno
- **Questions remained**
 - Were these SCP4 versions sufficient?
 - Might other restrictions of generalization be needed?
 - Had the SCP4 author to invent new modifications of SCP4?

Domain-Specific Special-Purpose MRSC

- **2010-2011 Andrei Klimov**

- have **theoretically** explained and proved that the **coverability problem** is solvable for **monotonic counter systems** by an **iterative procedure** of applying a **domain-specific supercompiler** for counter systems **varying a parameter** of generalization
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- **Algorithm**

- Scp_l – a domain-specific supercompiler for counter systems with parameter l **prohibiting generalization of integers** $n < l$
 - the **simplest** version: integers $n \geq l$ are **immediately generalized**
- **for** $l=1,2,3\dots$ **do**
 - use Scp_l to build a **residual set of configurations**
 - if** all residual configurations are **disjoint** with the target set $Unsafe$
 - then return** “Unreachable”
- This algorithm with the **simplest** supercompiler Scp_l **fits the ECC schema**
- My proof of its correctness **differs** from that of the ECC algorithmic schema
- ...and asserts a **stronger termination** statement:
 - it terminates for all **monotonic** counter systems and upper-closed *Unsafe sets*

ECC as a domain-specific multi-result supercompiler

- **2005-2007 Gilles Geeraerts *et al***
 - theory of ‘Expand, Enlarge and Check’ algorithmic schema (ECC) for solving the coverability problem of well-structured transition systems (WSTS)
 - *An MRSC for WSTS and its optimized versions*
- **Main ideas**
 - The set of all possible configurations \mathcal{C} is infinite (as usual)
 - *Def.* A finite $R \subseteq \mathcal{C}$ is called a residual set iff it is closed under “driving”:
 - $\text{Post}(\llbracket R \rrbracket) \subseteq \llbracket R \rrbracket$
 - Consider an ascending sequence of finite sets of configurations C_l :
 - $C_0 \subset C_1 \subset C_2 \subset C_3 \dots$
 - $\mathcal{C} = \bigcup C_l$
 - Consider residual sets $R \subseteq C_l$
 - The set $\{R \mid R \subseteq C_l\}$ of all such residual sets is finite as C_l is finite
 - Hence, it is solvable whether there exists a safe $R \subseteq C_l$ (that is, all configurations in R are disjoint with the target set *Unsafe*)
 - Iterate for $l=0,1,2,3\dots$
 - If a safe residual set $R \subseteq \mathcal{C}$ exists, then C_l s.t. $R \subseteq C_l$ exists and hence the iterative procedure terminates
- The notion of MRSC is wider

ECC as a domain-specific multi-result supercompiler

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- **Main ideas**

- The set of all possible configurations \mathcal{C}
- *Def.* A finite $R \subseteq \mathcal{C}$
 - $\text{Post}(\llbracket R \rrbracket) \subseteq R$
- Consider an ascending chain of residual sets $C_0 \subset C_1 \subset C_2 \subset \dots$
 - $\mathcal{C} = \bigcup C_l$
- Consider residual sets $R \subseteq \mathcal{C}$
- The set $\{R \mid R \subseteq \mathcal{C}\}$ is solvable (that is, all configurations in R are covered by some C_l)
- Iterate for $l=0,1,2,3,\dots$
- If a safe residual set R is found, then MRSC finds it and hence the iteration terminates

Where is the well-structuredness of a TS used?

(WS = monotonicity + well-quasi-order)

- Existence of safe R when the TS is safe
- Optimizations: reducing the search space

Without the well-structuredness:

If there exists an inductive proof that a TS is safe with the inductive hypotheses in form of a residual set of configurations, then MRSC finds it

- The notion of MRSC is wider

Related work: Supercompilation-like algorithms

1969 Richard M. Karp and Raymond E. Miller.

Parallel Program Schemata, *J. Comput. Syst. Sci.* 3(2), 147-195.

- a covering tree for ordinary Petri nets
- like supercompilation with lower-node generalization

1993 Alain Finkel.

The minimal coverability graph for Petri nets.

In Grzegorz Rozenberg, ed., *Advances in Petri Nets 1993*, LNCS 674, 210-243.

- an attempt to improve the KM algorithm
- generalization of upper configuration
- a flaw: under-approximation of the minimal coverability set

...and a lot of other works...

2006 Gilles Geeraerts, Jean-François Raskin, and Laurent Van Begin.

Expand, Enlarge and Check: New algorithms for the coverability problem of WSTS. *Journal of Computer and System Sciences*, 72(1), 180-203.